

Section 2.1

Suppose we want to approximate A using n rectangles.
Then we partition the interval $[a, b]$ into n subintervals
of width

$$\Delta x = \frac{b-a}{n}.$$

Then the first subinterval is $[a, a + \Delta x]$

second subinterval is $[a + \Delta x, a + 2\Delta x]$

:

i th subinterval is $[a + (i-1)\Delta x, a + i\Delta x]$.

:

n th subinterval is $[a + (n-1)\Delta x, a + n\Delta x]$

where $a + n\Delta x = a + n \cdot \frac{b-a}{n} = a + (b-a) = b$.

We call the right endpoint of the i th subinterval x_i

where

$$x_i = a + i\Delta x$$

Note that the left endpoint of the i th subinterval is,
therefore, x_{i-1} .

We can then denote $x_0 = a$ and $x_n = b$.

Now we assume that $f(x)$ is continuous on $[a, b]$ and therefore, on each subinterval. Then the Extreme Value Theorem guarantees that $f(x)$ will have both a minimum value and a maximum value on each subinterval.

Let $x = m_i$ be the absolute minimum of $f(x)$ on the i th subinterval, so $f(m_i)$ is the minimum value. Then we would choose the height of each rectangle to be $f(m_i)$ in order to compute the lower sum.

Similarly, let $x = M_i$ be the absolute maximum of $f(x)$ on the i th subinterval, so $f(M_i)$ is the maximum value. Then we choose $f(M_i)$ to be the height of the rectangles to obtain the upper sum.

In general, then, the lower sum is given by

$$\begin{aligned} s(n) &= f(m_1)\Delta x + f(m_2)\Delta x + f(m_3)\Delta x + \cdots + f(m_n)\Delta x \\ &= [f(m_1) + f(m_2) + f(m_3) + \cdots + f(m_n)]\Delta x \end{aligned}$$

The upper sum is given by

$$\begin{aligned} S(n) &= f(M_1)\Delta x + f(M_2)\Delta x + f(M_3)\Delta x + \cdots + f(M_n)\Delta x \\ &= [f(M_1) + f(M_2) + f(M_3) + \cdots + f(M_n)]\Delta x \end{aligned}$$

Consider a sum of n terms $a_1, a_2, a_3, \dots, a_n$ for which all of the terms can be described by a general formula a_i . Here we call i the index of summation.

Then we can write the sum in sigma notation as follows:

$$a_1 + a_2 + a_3 + \dots + a_n = \sum_{i=1}^n a_i$$

Here, we call a_i the summand, $i=1$ the lower bound of summation, and n the upper bound of summation.

eg $1 + 4 + 9 + 16 + \dots + 144$

All of the terms can be written in the form i^2 .

Thus $a_i = i^2$ with lower bound 1 and upper bound 12. Hence we can write this sum in

sigma notation as $\sum_{i=1}^{12} i^2$

eg $2 + \frac{3}{8} + \frac{4}{27} + \frac{5}{64} + \dots + \frac{n+1}{n^3} = \sum_{i=1}^n \frac{i+1}{i^3}$

Theorem : Basic Properties of Sums

$$\textcircled{1} \quad \sum_{i=1}^n k a_i = k \sum_{i=1}^n a_i \quad \text{for any constant } k$$

$$\textcircled{2} \quad \sum_{i=1}^n (a_i + b_i) = \sum_{i=1}^n a_i + \sum_{i=1}^n b_i$$

$$\textcircled{3} \quad \sum_{i=1}^n (a_i - b_i) = \sum_{i=1}^n a_i - \sum_{i=1}^n b_i$$

However, $\sum_{i=1}^n a_i b_i \neq \left(\sum_{i=1}^n a_i \right) \cdot \left(\sum_{i=1}^n b_i \right)$

$$\sum_{i=1}^n \frac{a_i}{b_i} \neq \frac{\sum_{i=1}^n a_i}{\sum_{i=1}^n b_i}$$

$$\begin{aligned} \text{eg } \sum_{i=1}^n (2i - 1) &= \sum_{i=1}^n 2i - \sum_{i=1}^n 1 \\ &= 2 \sum_{i=1}^n i - \sum_{i=1}^n 1 \end{aligned}$$

Theorem : Summation Formulas

$$\textcircled{1} \quad \sum_{i=1}^n 1 = \underbrace{| + | + | + | + \dots + |}_{n \text{ times}} = n$$

$$\textcircled{2} \quad \sum_{i=1}^n i = 1 + 2 + 3 + 4 + \dots + n = \frac{n(n+1)}{2}$$

$$\textcircled{3} \quad \sum_{i=1}^n i^2 = 1^2 + 2^2 + 3^2 + 4^2 + \dots + n^2 = \frac{n(n+1)(2n+1)}{6}$$

$$\textcircled{4} \quad \sum_{i=1}^n i^3 = 1^3 + 2^3 + 3^3 + 4^3 + \dots + n^3 = \frac{n^2(n+1)^2}{4}$$